

Exercises - C11, C++11

In the exercises below, you should use the CPPMEM tool, available online at:

<http://svr-pes20-cppmem.cl.cam.ac.uk/cppmem/>

1. Load CPPMEM, use the left-most drop down box to select the *examples/IRIW* test, change all memory orders to `memory_order_relaxed`, and press the run button. CPPMEM will print an execution graph of the relaxed behaviour that this test allows.
 - (a) Play with the tick boxes on the left to turn on and off the printing of relations. Make sure you look at *hb*, *rf*, *sc* and *mo*.
 - (b) Use the following table to “compile” the program to a Power program:

C++0x Operation	POWER Implementation
Non-atomic Load	<code>ld</code>
Load Relaxed	<code>ld</code>
Load Seq Cst	<code>sync; ld; cmp; bc; isync</code>
Non-atomic Store	<code>st</code>
Store Relaxed	<code>st</code>
Store Seq Cst	<code>sync; st</code>

Will the new program produce the relaxed behaviour on the Power abstract machine?

- (c) Recall that on Power, placing a `sync` barrier between each pair of reads was sufficient to forbid this behaviour. Press *reset*, and then alter the C program so that the last load on each load thread has order `memory_order_seq_cst`. How does the “compiled” program change, and is the relaxed behaviour still allowed on the Power abstract machine? Why is the relaxed behaviour still allowed in C?
 - (d) By choosing different memory orders for each load or store in the C program, forbid the relaxed behaviour in C. Why does your solution work?
2. Use the left-most drop down box to select the *examples/LB* test and press run.
 - (a) Could the SC machine reproduce the execution that CPPMEM prints?
 - (b) Press the *next consistent* button and inspect each consistent execution. Can all of the executions be seen on an SC machine?
 - (c) Use the *previous candidate* and *next candidate* buttons to explore the candidate executions. Find the one with the relaxed behaviour. What predicates in the model fail?
 - (d) Alter the memory orders in the program so that load-buffering relaxed behaviour can be witnessed.
3. For each program, enumerate the values that can be read at the commented line, given the commented constraints, and say why:

(a)

```

atomic_int x = 0;
x.store(1,relaxed);    | x.load(relaxed); \\ this reads 1
                       | x.load(relaxed); \\ What values can be read?

```

(b)

```

atomic_int x = 0;
x.store(1,relaxed);
x.load(relaxed); \\ What values can be read?
x.store(2,relaxed);

```