

$var, x$	term variable		
$tvar, X$	type variable		
$term, t$	$::=$		term
	$x$		variable
	$\lambda x . t$	bind $x$ in $t$	abstraction
	$t_1 t_2$		application
	$(t)$	S	
	$\{ t_1 / x \} t_2$	M	
$value, v$	$::=$		value
	$\lambda x . t$	bind $x$ in $t$	abstraction
$type, S, T$	$::=$		type
	$X$		variable
	$T \rightarrow T'$		function
	$(T)$	S	
$env, E$	$::=$		type environment
	<b>empty</b>		
	$env, x : T$		
$formula$	$::=$		
	<i>judgement</i>		
	$x : T \in env$		
	<b>ok</b> $env$		
$terminals$	$::=$		
	$\lambda$		
	$\longrightarrow$		
	$\Longrightarrow$		
	$\rightarrow$		
	$\vdash$		
	$\in$		
$Jtype$	$::=$		
	$E \vdash t : T$		
$Jop$	$::=$		
	$t_1 \longrightarrow t_2$		$t_1$ reduces to $t_2$
$Jopf$	$::=$		
	$t_1 \Longrightarrow t_2$		$t_1$ full-reduces to $t_2$
$judgement$	$::=$		
	<i>Jtype</i>		
	<i>Jop</i>		
	<i>Jopf</i>		
$user\_syntax$	$::=$		

| *var*  
 | *tvar*  
 | *term*  
 | *value*  
 | *type*  
 | *env*  
 | *formula*  
 | *terminals*

$E \vdash t : T$

$$\begin{array}{c}
 \mathbf{ok} E \\
 \frac{x : T \in E}{E \vdash x : T} \quad \text{TYPING\_VAR} \\
 \\
 \frac{E \vdash t_1 : S \rightarrow T \quad E \vdash t_2 : S}{E \vdash t_1 t_2 : T} \quad \text{TYPING\_APP} \\
 \\
 \frac{E, x : S \vdash t : T}{E \vdash \lambda x. t : S \rightarrow T} \quad \text{TYPING\_LAM}
 \end{array}$$

$t_1 \longrightarrow t_2$   $t_1$  reduces to  $t_2$

$$\begin{array}{c}
 \frac{}{(\lambda x. t_1) v_2 \longrightarrow \{v_2/x\} t_1} \quad \text{AX\_APP} \\
 \\
 \frac{t_1 \longrightarrow t'_1}{t_1 t_2 \longrightarrow t'_1 t_2} \quad \text{CTX\_APP\_FUN} \\
 \\
 \frac{t_2 \longrightarrow t'_2}{v_1 t_2 \longrightarrow v_1 t'_2} \quad \text{CTX\_APP\_ARG}
 \end{array}$$

$t_1 \Longrightarrow t_2$   $t_1$  full-reduces to  $t_2$

$$\begin{array}{c}
 \frac{}{(\lambda x. t_1) t_2 \Longrightarrow \{t_2/x\} t_1} \quad \text{FULL\_REDUCE\_RED} \\
 \\
 \frac{t_1 \Longrightarrow t'_1}{t_1 t_2 \Longrightarrow t'_1 t_2} \quad \text{FULL\_REDUCE\_APP1} \\
 \\
 \frac{t_2 \Longrightarrow t'_2}{t_1 t_2 \Longrightarrow t_1 t'_2} \quad \text{FULL\_REDUCE\_APP2} \\
 \\
 \frac{t_1 \Longrightarrow t'_1}{\lambda x. t_1 \Longrightarrow \lambda x. t'_1} \quad \text{FULL\_REDUCE\_ABS}
 \end{array}$$

Definition rules:            10 good      0 bad  
 Definition rule clauses: 20 good      0 bad